

(+ (100 50) 2)

Let's add local variables and **binary ops** to our compiler

```

expr := <number>
      | (add1 <expr>)
      | (let (<name> <expr>) <expr>)
      | <name>
      | (+ <expr> <expr>)

```

Result	Programs	Stack Layout	Assembly
---------------	-----------------	---------------------	-----------------

```
((+ 1 2) 3)
```

```
(+ 1 (+ 2 (+ 3 4)))
```

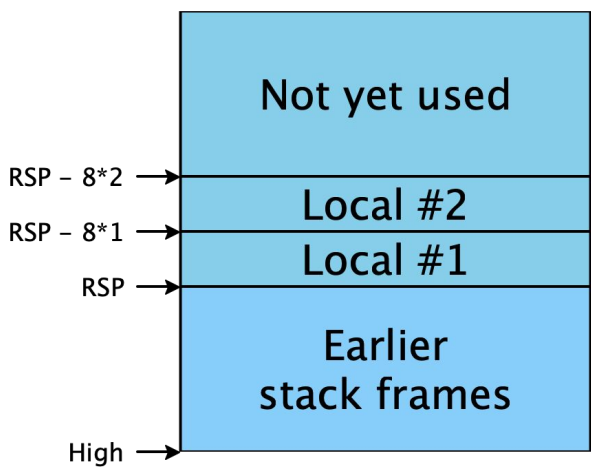
```
(let (x 10)
  (let (y 10)
    (+ x y)))
```

```
(let (x 10)
  (let (x (add1 x))
    (+ x 10)))
```

```
(+ (let (x 10) (add1 x))
  (let (y 7) (+ x y)))
```

```
(+ (let (x 10) (add1 x))
  (let (y 7) (+ 10 y)))
```

(+ e1 e2)



Let's add **local variables** and **binary ops** to our compiler

(+ (100 50) 2)

```
expr := <number>
      | (add1 <expr>)
      | (let (<name> <expr>) <expr>)
      | <name>
      | (+ <expr> <expr>)
```

What assembly is produced?

```
enum Expr {
  Num(i32),
  Add1(Box<Expr>),
}
```

```
enum Val {
  Reg(Reg),
  Imm(i32),
}
```

```
enum Reg {
  RAX,
}
```

```
enum Instr {
  IMov(Val, Val),
  IAdd(Val, Val),
  ISub(Val, Val),
  IMul(Val, Val),
}
```

```
fn compile(e:&Expr, env:&Env, is) -> String {
  match e {
    Expr::Num(n) => {
      format!("mov rax, {}", *n)
    }
    Expr::Add1(e1) => {
      let e1_code = compile(e1, env);
      format!("{e1_code}
              add rax, 1")
    }
    Expr::Var(x) => {
      let offset = env.get(x).unwrap();
      format!("mov rax, [rsp - 8*{offset}]")
    }
    Expr::Let(x, e1, e2) => {
      let e1_code = compile(e1, env, is);
      let new_env = env.update(x.clone(), offset);
      let e2_code = compile(e2, &new_env, is + 1);
      format!("{e1_code}
              mov [rsp - 8*{is}], rax
              {e2_code}")
    }
    Expr::Plus(e1, e2) => {
  }
}
}
```

What should these evaluate to? Why?

```
(let (x 5)
  (if (= x 10) (+ x 2) x))
```

```
(if 5 true false)
```

```
(+ 7 true)
```

```
(= true 1)
```

What should be in RAX after these are done evaluating? Why?

5

-3

true

false

(= 3 5)

(+ 4 7)

```
enum Expr {
  Num(i32),
  True, False,
  If(Box<Expr>, Box<Expr>, Box<Expr>)
  Eq(Box<Expr>, Box<Expr>)
  Add1(Box<Expr>),
  Plus(Box<Expr>, Box<Expr>),
  Let(String, Box<Expr>, Box<Expr>)
  Id(String),}

fn compile_expr(e : &Expr, si : i32, env : &HashMap<String, i32>) -> String {
  match e {
    Expr::Num(n) =>

    Expr::True =>

    Expr::False =>

    Expr::Add1(subexpr) => ...,
    Expr::Plus(e1, e2) => ...,
    Expr::Let(x, e, body) => ...,

    Expr::If(cond, thn, els) => {

  }

  Expr::Eq(e1, e2) => {

  }

}
}
```

This is 64 bits: 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000

This is 5: 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0101

This is 5 shifted 1 to the left, AKA 10: 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 1010

If we're OK with 63-bit numbers, can use LSB for tag 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 1010 = 5
0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0001 = false
0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0000 0011 = true

What does this mean for code generation?

What should we do the next time we need a new type? (string, heap-allocated object, etc.)

Condition Codes (that matter for us): Overflow, Sign, Zero
many instructions set these; arithmetic, shifting, etc. mov does not

cmp <reg>, <val> compute <reg> - <val> and set condition codes (value in <reg> does not change)
some cases to think about:

<reg> = -2^64, <val> = 1 Overflow: ___ Sign: ___ Zero: ___

<reg> = 0, <val> = 1 Overflow: ___ Sign: ___ Zero: ___

<reg> = 1, <val> = 0 Overflow: ___ Sign: ___ Zero: ___

<reg> = -1, <val> = -2 Overflow: ___ Sign: ___ Zero: ___

test <reg>, <val> perform bitwise and on the two values, but don't change <reg>, and set condition codes as appropriate. Useful for mask checking. test rax, 1 will set Z to true if and only if the LSB is 1

<label>: set this line as a label for jumping to later

jmp <label> unconditionally jump to <label>

jne <label> jump to <label> if Zero is not set (last cmped values not equal)

je <label> jump to <label> if Zero is set (last cmped values are equal)

jge <label> jump to <label> if Overflow is the same as Sign (which corresponds to >= for last cmp)

jle <label> jump to <label> if Zero set or Overflow != Sign (which corresponds to <= for last cmp)

shl <reg> shift <reg> to the left by 1, filling in least-significant bit with zero

sar <reg> shift <reg> to the right by 1, filling in most-significant bit to preserve sign

shr <reg> shift <reg> to the right by 1, filling in most-significant bit with zero